

### Andrew Suk, Csaba D. Tóth, Andres J. Ruiz-Vargas

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Outline	Thrackles 0000000	The main tool	Redrawing	Summary



- Thrackles
- Tangles
- Tangled-thrackles

## 2 The main tool

# 3 Redrawing





Outline	Thrackles ●000000	The main tool	Redrawing	Summary
Thrackles				
Thrackles				

• A drawing of a graph.



Outline	Thrackles ●oooooo	The main tool	Redrawing	Summary
Thrackles				
Thrackles				

- A drawing of a graph.
- Every pair of edges meets exactly once: at a vertex or at a crossing point.



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Outline	Thrackles oooooo	The main tool	Redrawing	Summary
Thrackles				
State of	f affairs			

- Conway's conjecture: If a thrackle has n vertices then it has at most n edges.
- If true, it would be tight: every cycle with more than 4 vertices can be drawn as a thrackle.



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- (2000)Cairns and Nikolayevsky: 1.5n
- (2011)Fulek, Pach: 1.428n



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Tangles				
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• Touching points are different.

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Tangles				
Tangles				

- A drawing of a graph.
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Outline	Thrackles ○○○●○○○	The main tool	Redrawing	Summary
Tangles				
Theorem				

### Theorem

(Pach, Tóth, Radoičić, 2011) A tangle with n vertices has at most n edges.

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Tangled-th	rackles			

• A drawing of a graph.



Outline	Thrackles ○○○○●○○	The main tool	Redrawing	Summary
Tangled-thrackles				
Tangled-th	rackles			

- A drawing of a graph.
- Every pair of edges meets exactly once: at a a vertex, at a crossing or at a touching.

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- A drawing of a graph.
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• Touching and crossing points are all distinct.

- A drawing of a graph.
- Every pair of edges meets exactly once: at a a vertex, at a crossing or at a touching.
- Touching and crossing points are all distinct.
- What is the maximum number of edges tangled-thrackle with *n* vertices can have? (Pach, Radoičić, and Tóth)

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Tangled-th	rackles			

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## • $O(n \log^{12} n)$ (Pach, Radoičić, and Tóth, 2012).

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- $O(n \log^{12} n)$  (Pach, Radoičić, and Tóth, 2012).
- Conjectured O(n).

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- $O(n \log^{12} n)$  (Pach, Radoičić, and Tóth, 2012).
- Conjectured *O*(*n*).
- We proved their conjecture.

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Tangled-thrackles				
A first obs	ervation: Bo	ounding numbe	r of touchings	5

• No 200 edges touch another set of 200 edges.



Tangled-thrackles

# A first observation: Bounding number of touchings

- No 200 edges touch another set of 200 edges.
- I.e. the touching graph has no K<sub>200,200</sub>. By Kövári, Sós, Turán number of touchings is at most

$$c|E|^{2-1/200} \le c(n\log^{12}n)^{2-1/200} \le cn^{2-1/1000}$$

Outline	Thrackles 0000000	The main tool	Redrawing	Summary

## Odd-crossing number

### Definition

The odd-cr(G) is the least number of pairs of edges that cross an odd number of times among all drawings of G.

### Definition

The bisection width b(G) is the least number of edges from  $V_1$  to  $V_2$  among all partitions  $V_1$ ,  $V_2$  of V with  $V_i \ge n/3$ .

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#### Theorem

(Pach, Tóth) There is an absolute constant  $c_2$  such that if G is a graph with n vertices of vertex degrees  $d_1, \ldots, d_n$ , then

$$b(G) \leq c_2 \log n \sqrt{\textit{odd-cr}(G) + \sum_{i=1}^n d_i^2}.$$

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### • How do we use this theorem?



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- How do we use this theorem?
- We assume *G* is bipartite. Whichever edges are touching we changed them slightly so that they become disjoint.

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• We show that if *G* is drawn as tangled thrackle then its odd-crossing number is small.



Figure : Redrawing procedure

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Outline	Thrackles 0000000	The main tool	Redrawing	Summary
Redrawing	)			

• After redrawing a pair of edges crosses an odd number of times if and only if they were originally touching.

# Bounding number of touchings

### Theorem

(Pach, Tóth) There is an absolute constant  $c_2$  such that if G is a graph with n vertices of vertex degrees  $d_1, \ldots, d_n$ , then

$$b(G) \leq c_2 \log n \sqrt{\textit{odd-cr}(G) + \sum_{i=1}^n d_i^2}.$$

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• odd-cr(G) = the number of touchings  $\leq cn^{2-1/1000}$ 

# Bounding number of touchings

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odd-cr(G) = the number of touchings ≤ cn<sup>2-1/1000</sup>
b(G) ≤ n<sup>1-1/2000</sup>

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Summary				

• Small bisection width: bound odd crossing number.



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• Redrawing and Kövári, Sós, Turán.

Outline	Thrackles 0000000	The main tool	Redrawing	Summary
Summary				

- Small bisection width: bound odd crossing number.
- Redrawing and Kövári, Sós, Turán.
- Decompose the graph into two parts using small bisection width and apply induction.
- Show that a tangled thrackle has at most  $c(n n^{1-1/4000})$  edges.

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Outline	Thrackles 0000000	The main tool	Redrawing	Summary
Some o	pen questio	ns		

• What is the smallest *t* such that there is no *K*<sub>*t*,*t*</sub> on the touching graph?

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• Thrackle conjecture is still open.

Outline	Thrackles 0000000	The main tool	Redrawing	Summary
Thank y	ou.			

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