A Schnyder-type drawing algorithm for 5-connected triangulations

Olivier Bernardi¹, Éric Fusy² and Shizhe Liang¹

- 1. Dept. of Math, Brandeis University
- 2. LIGM/CNRS, Université Gustave Eiffel

Graph drawing'23

Triangulation = graph embedded in the plane, all faces of degree 3



= embedded maximal planar graph

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- 3-connected \Leftrightarrow simple
- 4-connected \Leftrightarrow no separating 3-cycle
- 5-connected \Leftrightarrow no separating 4-cycle

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embedded maximal planar graph

4-connected not 5-connected

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3-connected case

Schnyder structures on simple triangulations [Schnyder'89]

Any triangulation admits a **labeling** of corners by $\{1, 2, 3\}$ satisfying



Schnyder structures on simple triangulations [Schnyder'89]

Yields 3 spanning trees T_1, T_2, T_3 (Schnyder wood)





inner vertex





tree-paths from v partition inner faces into 3 regions $R_1(v)$, $R_2(v)$, $R_3(v)$



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 $^{\mathsf{b}}v_3$







4-connected case

Triangulations of the 4-gon, irreducibility

A triangulation of the 4-gon is irreducible if all 3-cycles bound faces



not irreducible

irreducible

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Transversal structures

aka regular edge-labelings [He'93] (structures dual to rectangular tilings)



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yields two bipolar orientations:





















9 faces in blue map



















leftmost outgoing red path

































leftmost outgoing blue path











leftmost outgoing blue path + rightmost ingoing blue path











leftmost outgoing blue path + rightmost ingoing blue path











leftmost outgoing blue path + rightmost ingoing blue path























planar straight-line drawing











planar straight-line drawing



Face-counting algorithm on square grid v_2






Face-counting algorithm on square grid v_2









 18×18 grid

Face-counting algorithm on square grid









Face-counting algorithm on square grid v_3









Face-counting algorithm on square grid v_2









4-wood associated to transversal structure



left outgoing blue edges

left outgoing red edges

4-wood associated to transversal structure



4-wood associated to transversal structure



















































5-connected case

5c-triangulations

5c-triangulation = triangulation of 5-gon such that every cycle with at least one vertex inside has length ≥ 5



not 5c-triangulation



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not 5c-triangulation



5c-triangulation $\Uparrow \approx$

triangulation augmented by v_∞ is 5-connected

5c-labelings

[Bernardi, F, Liang'23]

Any 5c-triangulation has a labeling of corners by $\{1, 2, 3, 4, 5\}$ so that































• Linear time complexity

- Linear time complexity
- Displays rotational symmetries



A:
$$(2,6,4,2,1)$$

B: $(1,2,6,4,2)$
C: $(2,1,2,6,4)$
D: $(4,2,1,2,6)$
E: $(6,4,2,1,2)$
F: $(3,3,3,3,3)$



- Linear time complexity
- Displays rotational symmetries



• Variations: weighted faces, vertex-counting

- Linear time complexity
- Displays rotational symmetries

- Variations: weighted faces, vertex-counting
- Vertex resolution better than in the 3- or 4-connected drawings smallest distance between vertices

(drawing normalized to have outer k-gon inscribed in circle of radius 1)

new proof of existence for 4-connected (from 3-connected)



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new proof of existence for 4-connected (from 3-connected)



similar proof of existence for 5-connected (from 4-connected)