# A Schnyder-type drawing algorithm for 5-connected triangulations 

Olivier Bernardi ${ }^{1}$, Éric Fusy ${ }^{2}$ and Shizhe Liang ${ }^{1}$

1. Dept. of Math, Brandeis University
2. LIGM/CNRS, Université Gustave Eiffel

Graph drawing'23

## Triangulations

Triangulation $=$ graph embedded in the plane, all faces of degree 3

$=$ embedded maximal planar graph

## Triangulations

Triangulation $=$ graph embedded in the plane, all faces of degree 3

$=$ embedded maximal planar graph

For any graph (def):
$k$-connected: deleting any subset of $<k$ vertices does not disconnect

## Triangulations

Triangulation $=$ graph embedded in the plane, all faces of degree 3

$=$ embedded maximal planar graph

For any graph (def):
$k$-connected: deleting any subset of $<k$ vertices does not disconnect

For triangulations:

| 3-connected | $\Leftrightarrow$ | simple |
| :--- | :--- | :---: |
| 4-connected | $\Leftrightarrow$ | no separating 3-cycle |
| 5-connected | $\Leftrightarrow$ | no separating 4-cycle |

## Triangulations

Triangulation $=$ graph embedded in the plane, all faces of degree 3

$=$ embedded maximal planar graph

4-connected

## not 5-connected

For any graph (def):
$k$-connected: deleting any subset of $<k$ vertices does not disconnect

For triangulations:

| 3-connected | $\Leftrightarrow$ | simple |
| :--- | :--- | :---: |
| 4-connected | $\Leftrightarrow$ | no separating 3-cycle |
| 5-connected | $\Leftrightarrow$ | no separating 4-cycle |

## 3-connected case

Schnyder structures on simple triangulations [Schnyder'89]

Any triangulation admits a labeling of corners by $\{1,2,3\}$ satisfying


Schnyder structures on simple triangulations [Schnyder'89]

Yields 3 spanning trees $T_{1}, T_{2}, T_{3} \quad$ (Schnyder wood)


inner vertex


Schnyder's face-counting algorithm
[Schnyder'90]

[Schnyder'90] tree-paths from $v$ partition inner faces into 3 regions $R_{1}(v), R_{2}(v), R_{3}(v)$

[Schnyder'90]
tree-paths from $v$ partition inner faces into
3 regions $R_{1}(v), R_{2}(v), R_{3}(v)$


## Schnyder's face-counting algorithm

tree-paths from $v$ partition inner faces into
3 regions $R_{1}(v), R_{2}(v), R_{3}(v)$


Schnyder's face-counting algorithm
[Schnyder'90]

planar straight-line drawing


## Schnyder's face-counting algorithm

[Schnyder'90]
planar straight-line drawing


## Schnyder's face-counting algorithm

[Schnyder'90]
planar straight-line drawing


## 4-connected case

## Triangulations of the 4-gon, irreducibility

A triangulation of the 4-gon is irreducible if all 3-cycles bound faces


## Triangulations of the 4-gon, irreducibility

A triangulation of the 4-gon is irreducible if all 3-cycles bound faces

not irreducible


I
triangulation augmented by $v_{\infty}$ is 4-connected
aka regular edge-labelings (structures dual to rectangular tilings)


A 4-triangulation admits a transversal structure iff it is irreducible
aka regular edge-labelings (structures dual to rectangular tilings)


A 4-triangulation admits a transversal structure iff it is irreducible
yields two bipolar orientations:


Face-counting algorithm


Face-counting algorithm


Face-counting algorithm

$10 \times 9$ grid

Face-counting algorithm


Face-counting algorithm



Face-counting algorithm



Face-counting algorithm



Face-counting algorithm



Face-counting algorithm



leftmost outgoing blue path + rightmost ingoing blue path


Face-counting algorithm

leftmost outgoing blue path + rightmost ingoing blue path

Face-counting algorithm

leftmost outgoing blue path + rightmost ingoing blue path


Face-counting algorithm


Face-counting algorithm

planar straight-line drawing

Face-counting algorithm

cone property (implies planarity)

planar straight-line drawing


Face-counting algorithm on square grid (


Face-counting algorithm on square grid (8 inner faces

$18 \times 18$ grid

Face-counting algorithm on square grid



Face-counting algorithm on square grid




## 4-wood associated to transversal structure



## 4-wood associated to transversal structure

yields 4 regions for each vertex $v$

right incoming red edges

$T_{4}$

left outgoing blue edges
right incoming blue edges

left outgoing red edges

## 4-wood associated to transversal structure

yields 4 regions for each vertex $v$

square-grid algo I
barycentric placement
right incoming red edges

$T_{4}$

left outgoing blue edges
right incoming blue edges

left outgoing red edges
(place $v$ at $\frac{4}{28} v_{1}+\frac{8}{28} v_{2}+\frac{4}{28} v_{3}+\frac{2}{28} v_{4}$ )

4-labeling associated to transversal structure


4-labeling associated to transversal structure


4-labeling associated to transversal structure


4-labeling associated to transversal structure


4-labeling associated to transversal structure


4-labeling associated to transversal structure


## 5-connected case

## 5c-triangulations

$5 c$-triangulation $=$ triangulation of 5 -gon such that every cycle with at least one vertex inside has length $\geq 5$

not 5c-triangulation


5c-triangulation

## 5c-triangulations

$5 c$-triangulation $=$ triangulation of 5 -gon such that every cycle with at least one vertex inside has length $\geq 5$

not 5c-triangulation


5c-triangulation

$$
\Uparrow \approx
$$

triangulation augmented by $v_{\infty}$
is 5 -connected

## 5c-labelings

Any 5 c -triangulation has a labeling of corners by $\{1,2,3,4,5\}$ so that

outer vertices

inner vertices


5c-labeling


$T_{3}$



$T_{3}$


configuration at inner vertex
[Felsner, Schrezenmaier, Steiner'20] other 5-woods (less restrictive) associated to pentagon-contact representations

$T_{3}$


Face-counting algorithm
[Bernardi,F,Liang'23]


Face-counting algorithm
[Bernardi,F,Liang'23]


Face-counting algorithm


## Face-counting algorithm



45 inner faces in total

place $v$ at

$$
\frac{4}{45} v_{1}+\frac{4}{45} v_{2}+\frac{18}{45} v_{3}+\frac{7}{45} v_{4}+\frac{12}{45} v_{5}
$$

Face-counting algorithm
[Bernardi,F,Liang'23]


## Face-counting algorithm


cone property



## Face-counting algorithm


cone property

(s)



5


Face-counting algorithm

$v_{1} \quad v_{5}$


RE: Not a grid drawing
cone property

$$
\notin
$$


sheer

coordinates in $\mathbb{Q}(\sqrt{5})$

Properties and variations

- Linear time complexity


## Properties and variations

- Linear time complexity
- Displays rotational symmetries


A: $(2,6,4,2,1)$
B: $(1,2,6,4,2)$
C: $(2,1,2,6,4)$
D: $(4,2,1,2,6)$
E: $(6,4,2,1,2)$
F: $(3,3,3,3,3)$


## Properties and variations

- Linear time complexity
- Displays rotational symmetries


$$
\begin{aligned}
& \text { A: }(2,6,4,2,1) \\
& \text { B: }(1,2,6,4,2) \\
& \text { C: }(2,1,2,6,4) \\
& \text { D: }(4,2,1,2,6) \\
& \text { E: }(6,4,2,1,2) \\
& \text { F: }(3,3,3,3,3)
\end{aligned}
$$



- Variations: weighted faces, vertex-counting


## Properties and variations

- Linear time complexity
- Displays rotational symmetries


A: $(2,6,4,2,1)$
B: $(1,2,6,4,2)$
C: $(2,1,2,6,4)$
D: $(4,2,1,2,6)$
E: $(6,4,2,1,2)$
F: $(3,3,3,3,3)$


- Variations: weighted faces, vertex-counting
- Vertex resolution better than in the 3- or 4-connected drawings


## smallest distance between vertices

(drawing normalized to have outer $k$-gon inscribed in circle of radius 1 )

## Strategy for proof of existence

new proof of existence for 4-connected (from 3-connected)


## Strategy for proof of existence

new proof of existence for 4-connected (from 3-connected)


## Strategy for proof of existence

 new proof of existence for 4-connected (from 3-connected)
bicolored
corner
unicolored corner

## Strategy for proof of existence

 new proof of existence for 4-connected (from 3-connected)
new proof of existence for 4-connected (from 3-connected)

outdegree 4
outdegree 1

## Strategy for proof of existence

 new proof of existence for 4-connected (from 3-connected)
outdegree 4
outdegree 1
no separating triangle I
orientation is "co-accessible" ( $\exists$ co-accessibility spanning tree)


Schnyder orientation

## outdegre

## Strategy for proof of existence

 new proof of existence for 4-connected (from 3-connected)
outdegree 4
outdegree 1



## Strategy for proof of existence

 new proof of existence for 4-connected (from 3-connected)
outdegree 4
outdegree 1

outdegree 3
similar proof of existence for 5-connected (from 4-connected)

